## Harold's Cryptology Cheat Sheet

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## Definitions

| Term |  |
| :--- | :--- |
| Cryptology | The study of cryptography and cryptoanalysis |
| Cryptography | Methods of encipherment (secret techniques) |
| Cryptoanalysis | Methods of decipherment (code breaking) |
| Plain | Plain text message to be encrypted |
| Cipher | Encrypted text message to be decrypted |
| Key | Secret string or set of numbers used to encrypt plain text |
| Stegenography | Information hiding in files, like JPG images |

## Text to Numbers Encoding

| Letter | Number | Letter | Number |
| :---: | :---: | :---: | :---: |
| $\mathbf{A}$ | 00 | $\mathbf{N}$ | 13 |
| $\mathbf{B}$ | 01 | $\mathbf{O}$ | 14 |
| $\mathbf{C}$ | 02 | $\mathbf{P}$ | 15 |
| $\mathbf{D}$ | 03 | $\mathbf{Q}$ | 16 |
| $\mathbf{E}$ | 04 | $\mathbf{R}$ | 17 |
| $\mathbf{F}$ | 05 | $\mathbf{S}$ | 18 |
| $\mathbf{G}$ | 06 | $\mathbf{T}$ | 19 |
| $\mathbf{H}$ | 07 | $\mathbf{U}$ | 20 |
| $\mathbf{I}$ | 08 | $\mathbf{V}$ | 21 |
| $\mathbf{J}$ | 09 | $\mathbf{W}$ | 22 |
| $\mathbf{K}$ | 10 | $\mathbf{X}$ | 23 |
| $\mathbf{L}$ | 11 | $\mathbf{Y}$ | 24 |
| $\mathbf{M}$ | 12 | $\mathbf{Z}$ | 25 |
|  |  | <space> | 26 |

## Cipher Methods

| Method | Concept | Example | How |
| :---: | :---: | :--- | :--- |

## Spreadsheet Example - Caesar Cipher

## Function Description Excel Formula

| Operation |  | A | B | C | $\mathbf{D}$ | $\mathbf{E}$ | F | $\mathbf{G}$ | $\mathbf{H}$ | $\mathbf{I}$ | J |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Plain Text | $\mathbf{1}$ | S | K | Y | I | S | C | L | E | A | R |
| Plain Text as \# | $\mathbf{2}$ | 18 | 10 | 24 | 8 | 18 | 2 | 11 | 4 | 0 | 17 |
| Cipher Text as \# | $\mathbf{3}$ | 25 | 27 | 5 | 15 | 25 | 9 | 18 | 11 | 7 | 24 |
| Cipher Text | $\mathbf{4}$ | Z | R | F | P | Z | J | S | L | H | Y |


| CODE("A") | Converts an ASCII character into a number | A2=CODE (A1) - CODE ("A") |
| :---: | :---: | :---: |
| $\operatorname{MOD}(\mathrm{n}, \mathrm{m})$ | Adds a fixed offset to each number ( n ) then mods it by $m$ | A3=MOD (A2 + 7, 26) |
| CHAR(65) | Converts a number into an ASCII character | A4 $=\operatorname{CHAR}\left(\mathrm{A} 3+\operatorname{CODE}\right.$ ( ${ }^{\text {a }}$ ") ) |
| Combined | All three functions combined into one | ```A4=CHAR (MOD (CODE (A1) - CODE("A") + 7, 26) + CODE("A"))``` |

## Frequency Analysis



## RSA Algorithm

| Term | Definition |
| :---: | :---: |
| RSA | Public key cryptosystem developed by Rivest, Adelman, and Shamir in 1978. |
| Key Prep | 1. Bob selects two large prime numbers, $p$ and $q$. <br> 2. Bob computes $\mathrm{N}=\mathrm{pq}$ and $\phi=(\mathrm{p}-1)(\mathrm{q}-1)$ <br> 3. Bob finds an integer e such that $\operatorname{gcd}(e, \phi)=1$. <br> 4. Bob computes the multiplicative inverse of $e \bmod \phi:$ an integer $d$ such that $(e d \bmod \phi)=1$. <br> 5. Public (encryption) key: N and e. <br> 6. Private (decryption) key: d. |
| Example | 1. Bob selects two primes: $\begin{aligned} & p=31 \\ & q=59 \end{aligned}$ <br> 2. Compute: $\begin{aligned} & N=p \cdot q=31 \cdot 59=1829 \\ & \phi=(p-1) \cdot(q-1)=30 \cdot 58=1740 \end{aligned}$ <br> 3. Find integer e such that gcd $(\mathrm{e}, \phi)=1$ <br> Guess e $=859$ and check: $\operatorname{gcd}(859,1740)=1$ <br> If the first guess is not relatively prime to $\phi$, try another. <br> 4. Using Euclid's algorithm, find A and B such that A $859+B \cdot 1740=1$ $\begin{aligned} & 79 \cdot 859+(-39) \cdot 1740=1 \\ & 79 \cdot 859=1 \bmod 1740 \\ & d=79 \text { is the inverse of } 859 \bmod 1740 \end{aligned}$ <br> 5. Public key: $(\mathrm{e}, \mathrm{N})$ $\begin{aligned} & e=859 \\ & N=1829 \end{aligned}$ <br> 6. Private key: (d, N) $\begin{aligned} & d=79 \\ & N=1829 \end{aligned}$ |
| Encryption | c $=m^{e} \bmod N \times$ |
| Decryption | $m=c^{d} \bmod N$ Private key: |
| Number Theory Fact | Let p and q be prime numbers and $\mathrm{pq}=\mathrm{N}$. Suppose that $\mathrm{m} \in \mathbf{Z}_{\mathrm{N}}$ and $\operatorname{gcd}(\mathrm{m}, \mathrm{N})=1$. Then $\mathrm{m}^{(p-1)(q-1)} \bmod \mathrm{N}=1$. |
| Theorem: Validity of the RSA Cryptosystem | If $m \in Z_{N}$ and $g c d(m, N)=1$, then RSA encryption and decryption applied to $m$ always yield m as the unique result. |

## Sources:

- SNHU MAT 230 - Discrete Mathematics, zyBooks.
- SNHU MAT 260 - Cryptology, Invitation to Cryptology, $1^{\text {st }}$ Edition, Thomas Barr, 2001.

